## An example of a proof by structural induction

Define the set  $\mathcal{B}$  of *binary trees* as follows:

- 1. A tree with a single node r is in  $\mathcal{B}$ ; and
- 2. If r is a node and  $T_1$  and  $T_2$  are binary trees, i.e.,  $T_1 \in \mathcal{B}$  and  $T_2 \in \mathcal{B}$ , then the tree  $T = (r, T_1, T_2)$  is a binary tree, i.e., T is in  $\mathcal{B}$ . You should view T as a tree with root r with r having as left child the tree  $T_1$  and as right child the tree  $T_2$ .

We now define two functions  $|.| : \mathcal{B} \to \mathbf{N}$  and  $h : \mathcal{B} \to \mathbf{N}$  which respectively return the number of nodes in a tree and the height of the tree.

- 1. Definition of  $|.|: \mathcal{B} \to \mathbf{N}$ :
  - (a) |r| = 1
  - (b)  $|(r, T_1, T_2)| = 1 + |T_1| + |T_2|$
- 2. Definition of  $h : \mathcal{B} \to \mathbf{N}$ :
  - (a) h(r) = 0(b)  $h((r, T_1, T_2)) = 1 + \max(h(T_1), h(T_2)).$

We will now proof that  $\forall T \in \mathcal{B}: |T| \leq 2^{h(T)+1} - 1.$ 

**Proof**. The proof will be by structural induction.

- 1. Let T be the tree consisting of a single node r. This is the base case. By definition, |r| = 1. By definition h(r) = 0. So clearly,  $|r| \le 2^{h(r)+1} 1$ .
- 2. Now let  $T = (r, T_1, T_2)$ . Since  $T_1$  and  $T_2$  are "simpler" trees than T, we can assume by *structural induction* that

$$|T_1| \le 2^{h(T_1)+1} - 1 \qquad (a)$$

and

$$|T_2| \le 2^{h(T_2)+1} - 1$$
 (b).

Therefore,

$$\begin{aligned} |T| &= |(r, T_1, T_2)| \\ &= 1 + |T_1| + |T_2| \quad \text{(by definition of } |.|) \\ &\leq 1 + (2^{h(T_1)+1} - 1) + (2^{h(T_2)+1} - 1) \quad \text{(by (a) and (b) above)} \\ &\leq 2.2^{\max(h(T_1),h(T_2))+1} - 1 \quad \text{(by simple algebra of max function)} \\ &= 2.2^{h(T)} - 1 \quad \text{(by definition of the } h \text{ function)} \\ &= 2^{h(T)+1} - 1. \end{aligned}$$